

MIXED-TYPE SECRETARY PROBLEMS ON SEQUENCES  
OF BIVARIATE RANDOM VARIABLES

01200424

坂口 実 (MINORU SAKAGUCHI)

Tech. U. Wrocław K. Szajowski (KRZYSZTOF SZAJOWSKI)

**ABSTRACT.** An employer interviews a finite number  $n$  of applicants for a position. They are interviewed one by one sequentially in random order. As each applicant  $i$  is interviewed, two attributes are evaluated by the amounts  $X_i$  and  $Y_i$ , where  $X_i$  may be "talent" (or quality), and  $Y_i$  may be the "look" (or degree of favorable impression) of the applicant. Suppose that  $\{X_i\}_{i=1}^n, \{Y_i\}_{i=1}^n$  is under the condition of full (no) -information secretary problem and that  $X_i$ 's and  $Y_i$ 's are mutually independent. We consider the three kinds of the employer's object and for each of three cases the problem is formulated by dynamic programming, and the optimal policy is explicitly derived.

1. INTRODUCTION (略)

2. SELECTING GOOD QUALITY UNDER THE REQUIREMENT OF THE BEST LOOK

Let  $\{(X_i, Y_i)\}_{i=1}^n$  be a sequence of independent bivariate r.v.s as given in the previous section. Observing the sequence  $(X_i, Y_i), i = 1, 2, \dots, n$ , one by one sequentially, we want to maximize  $EX_\tau$ , where  $\tau$  is the stopping time, under the required condition that  $Y_\tau = 1$  and  $2 \leq Y_t \leq t$ , for  $\tau < t \leq n$ .

We define state  $(i, x|n)$  which means that: no stop has yet been made, and we face the  $i$ -th object whose second attribute is relatively best (i.e.  $Y_i = 1$ ), and the first attribute is found as  $X_i = x$ .

Denoting, by  $V_{n,i}(x)$ , the expected reward obtained by employing the optimal stopping rule for the  $n$ -object problem at state  $(i, x|n)$ , we easily have

$$v_{n,i}(x) = \max\left\{\frac{i}{n}x, \sum_{j=i+1}^n \frac{i}{j(j-1)} E_F v_{n,j}(X)\right\}, \quad (2.1)$$

$i = 1, 2, \dots, n; v_{n,n}(x) \equiv x$ .

**Theorem 1.** The optimal stopping rule for the optimal equation (2.1) is: "Stop at the earliest object  $(X_i, Y_i)$  whose relative rank  $Y_i$ , when it appears, is unity and satisfies  $X_i > d_{n,i}$ ".

The sequence  $\{d_{n,i}\}_{i=1}^n$  is determined by the recursion

$$d_{n,i} = \frac{S_F(d_{n,i+1})}{i} + d_{n,i+1}, \quad (2.3)$$

( $i = 1, 2, \dots, n-1; d_{n,n} = 0$ ) where  $S_F(t) \equiv E_F(X \vee t)$ .

The optimal expected reward is given by  $E_F v_{n,1}(X)$ , i.e.  $\frac{S_F(d_{n,1})}{n}$ .

3. SELECTING GOOD QUALITY UNDER THE REQUIREMENT OF ONE OF THE TWO BEST IN THE LOOK. (略)

4. SELECTING GOOD LOOK UNDER THE REQUIREMENT OF THE BEST QUALITY

Let  $\{X_i\}_{i=1}^n$  be an iid sequence of r.v.s obeying uniform distribution over the unit interval  $0 \leq x \leq 1$ . Observing the sequence  $\{(X_i, Y_i)\}_{i=1}^n$ , one by one sequentially, we want to maximize  $\sum_{r=y}^n P(R_\tau = r | Y_\tau = y) a_r$ , where  $\tau$  is the stopping time, and  $a_r$  is the reward obtained when the absolute rank  $R_\tau$  of the second attribute at time  $\tau$ , is  $r$ , under the required

condition that  $X_\tau = \max_{1 \leq i \leq n} X_i$ . So the problem is, by following the notation in the previous section,

$$E\{a_{R_r(Y_1, \dots, Y_n)}; X_\tau = \max_{1 \leq i \leq n} X_i\} \rightarrow \max_r.$$

We assume that  $a_1 \geq a_2 \geq \dots \geq a_n \geq 0$ .

Define state  $(x, y_i | n)$  to mean that no stop has yet been made, and we face the  $i$ -th object  $(x_i, Y_i)$  with  $Y_i = y_i$  and  $X_i = \max(X_1, \dots, X_i) = x$ .

Denoting, by  $v_{n,i}(x, y)$ , the expected reward obtained by employing the optimal stopping rule for the  $n$ -period problem at state  $(x, y_i | n)$ , we easily have

$$v_{n,i}(x, y_i) = \max \left[ \frac{x^{n-i}}{\binom{n}{i}} \sum_{r=y_i}^n \binom{r-1}{y_i-1} \binom{n-r}{i-y_i} a_r, \sum_{j=i+1}^n \frac{x^{j-i-1}}{j} \sum_{w=1}^j \int_x^1 v_{n,j}(z, w) dz \right], \quad (4.2)$$

where  $1 \leq y_i \leq i$ ,  $1 \leq i \leq n$ ,  $0 \leq x \leq 1$ , with  $v_{n,n}(x, y) = a_y$ .

We discuss the following two simple cases.

Case 1.  $a_1 = 1, a_2 = a_3 = \dots = a_n = 0$ .

Let  $d_{n,i}$  ( $i = 1, 2, \dots, n-1$ ;  $d_{n,n-1} = n^{-1}$ ,  $d_{n,n} \equiv 0$ ) be a unique root in  $[0, 1]$  of the equation

$$\frac{1}{i} \sum_{l=1}^{n-i} \frac{x^{-l} - 1}{l} = 1. \quad (4.6)$$

For some small  $n$ , we find that  $d_{2,1} = 1/2$ ,  $d_{3,2} = 1/3$ ,  $d_{3,1} = \frac{1+\sqrt{6}}{5} \approx 0.6899$ ,  $d_{4,i} = 1/4$ ,  $\frac{i+\sqrt{8}}{7}$  ( $\approx 0.5469$ ),  $0.7755$  for  $i = 3, 2, 1$  respectively.  $d_{5,i} = 1/5$ ,  $\frac{1+\sqrt{10}}{9}$  ( $\approx 0.4625$ ),  $0.6591$ ,  $0.8248$  for  $i = 4, 3, 2, 1$  respectively and so on.

**Theorem 3.** The optimal stopping region for the optimality equation (4.2) in Case 1, is: "stop at the earliest  $(X_i, Y_i)$  that satisfies  $Y_i = 1$  and  $X_i = \max_{1 \leq t \leq i} X_t > d_{n,i}$ , where each  $d_{n,i}$  is given by a unique root of the equation 4.6.

Case 2.  $a_1 = a_2 = 1, a_3 = a_4 = \dots = a_n = 0$ .

For this case the optimality equation (4.2) becomes

$$\begin{aligned} v_{n,i}(x) &= \max \left\{ \frac{2n-i-1}{n(n-1)} x^{n-i}, \Psi_{n,i}(x) \right\} \quad \text{for states } (x, y_i = 1 | n) \\ u_{n,i}(x) &= \max \left\{ \frac{i(i-1)}{n(n-1)} x^{n-i}, \Psi_{n,i}(x) \right\} \quad \text{for states } (x, y_i = 2 | n). \end{aligned} \quad (4.10)$$

where

$$\Psi_{n,i}(x) \equiv \sum_{j=i+1}^n \frac{x^{j-i-1}}{j} \int_x^1 (v_{n,j}(z) + u_{n,j}(z)) dz,$$

**Theorem 4.** (田名) 文献:

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